204320 - Database Management

Chapter 5

More SQL: Complex Queries, Triggers, Views, and Schema Modification

Adapted for 204320

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Chapter 5 Outline

- More Complex SQL Retrieval Queries
- More Complex SQL Insert, Update, and Delete
- Specifying Constraints as Assertions and Actions as Triggers
- Views (Virtual Tables) in SQL
- Schema Change Statements in SQL

More Complex SQL Retrieval Queries

- Additional features allow users to specify more complex retrievals from database:
 - Nested queries
 - Joined tables
 - Outer joins
 - Aggregate functions and grouping

Comparisons Involving NULL and Three-Valued Logic

- Meanings of NULL
 - Unknown value: ไม่รู้ค่า
 - Unavailable or withheld value: ไม่ได้บันทึกค่า
 - Not applicable attribute: ไม่เกี่ยวข้องเลยไม่มีค่า
- Each individual NULL value considered to be different from every other NULL value
- SQL uses a three-valued logic:
 - TRUE, FALSE, and UNKNOWN

Comparisons Involving NULL and Three-Valued Logic (cont'd.)

 Table 5.1
 Logical Connectives in Three-Valued Logic

(a)	AND	TRUE	FALSE	UNKNOWN
	TRUE	TRUE	FALSE	UNKNOWN
	FALSE	FALSE	FALSE	FALSE
	UNKNOWN	UNKNOWN	FALSE	UNKNOWN
(b)	OR	TRUE	FALSE	UNKNOWN
	TRUE	TRUE	TRUE	TRUE
	FALSE	TRUE	FALSE	UNKNOWN
	UNKNOWN	TRUE	UNKNOWN	UNKNOWN
(c)	NOT			
	TRUE	FALSE		
	FALSE	TRUE		
	UNKNOWN	UNKNOWN		

Comparisons Involving NULL and Three-Valued Logic (cont'd.)

• SQL allows queries that check whether an attribute value is NULL: สามารถใช้คำสั่ง ธQL ตรวจสอบ ค่าที่เก็บอยู่ใน attribute ว่าเป็น NULL หรือไม่ โดยใช้คำสั่งดังแสดงด้านล่าง

- IS or IS NOT NULL

Query 18. Retrieve the names of all employees who do not have supervisors.

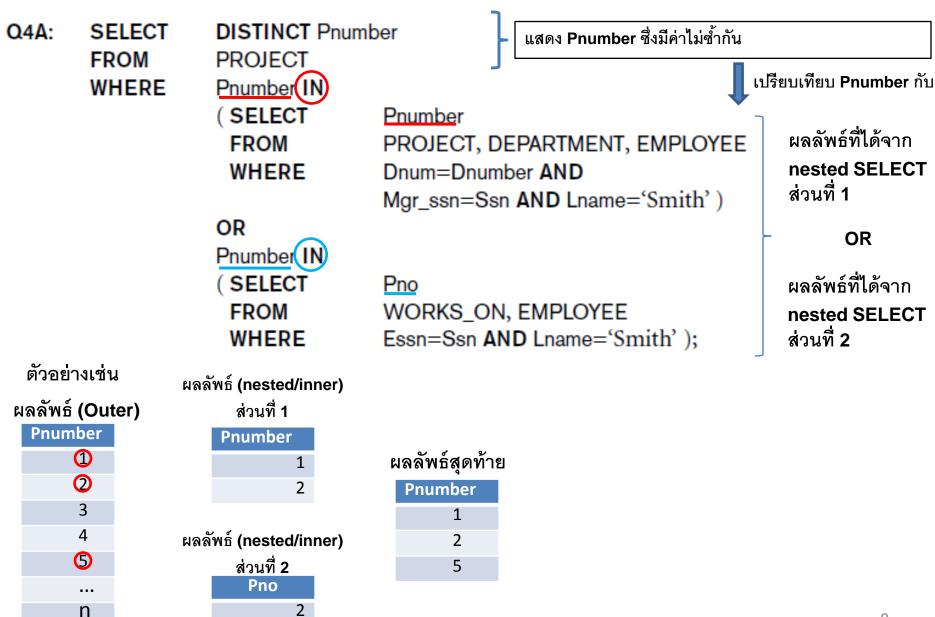
Q18: SELECT Fname, Lname

FROM EMPLOYEE

WHERE Super_ssn IS NULL;

Nested Queries, Tuples, and Set/Multiset Comparisons

- Nested queries: ประกอบด้วย 2 ส่วน คือ query ส่วนที่อยู่ ด้านใน และ query ส่วนที่อยู่ด้านนอก
 - Complete select-from-where blocks within WHERE clause of another query: ส่วนที่อยู่ด้านใน คือ select-from-where ที่อยู่หลัง keyword WHERE
 - Outer query: query ส่วนที่อยู่ด้านนอก
- Comparison operator IN: operator ที่ใช้ เปรียบเทียบระหว่าง outer และ inner
 - Compares value v with a set (or multiset) of values V
 - Evaluates to TRUE if v is one of the elements in V_{γ}



- Use tuples of values in comparisons
 - Place them within parentheses: หากต้องการเปรียบเทียบ

```
มากกว่า 1 attribute ให้ใส่วงเล็บ
```

SELECT DISTINCT Essn
FROM WORKS_ON

WHERE (Pno, Hours) IN (SELECT

FROM

WHERE

หมายเหตุ: ตอน Select มาทุก field (รวมทั้ง Pno และ Hours) แต่ผลลัพธ์สุดท้ายให้แสดงเฉพาะ field ที่ระบุหลัง keyword SELECT

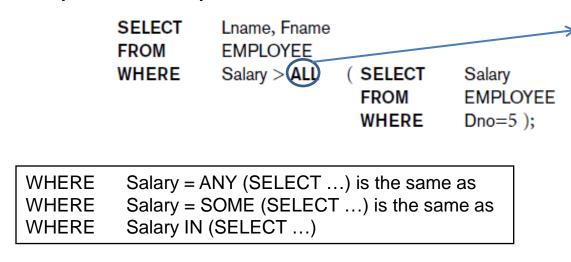
Pno, Hours
WORKS_ON
Essn='123456789');

The following SQL statement selects all customers with a City of "Paris" or "London ".

```
SELECT * // แสดงทุก fields ของ Customers
FROM Customers
WHERE City IN ('Paris','London'); // เปรียบเทียบ field City ที่ได้กับ Paris และ London
```

 Use other comparison operators to compare a single value v

- = ANY (or = SOME) operator
 - Returns TRUE if the value v is equal to some value in the set V and is hence equivalent to IN
- Other operators that can be combined with ANY (or SOME): >, >=, <, <=, and <>



ผลลัพธ์คือ พนักงานที่มีเงินเดือน มากกว่าทุกคนที่อยู่ในแผนก หมายเลข 5

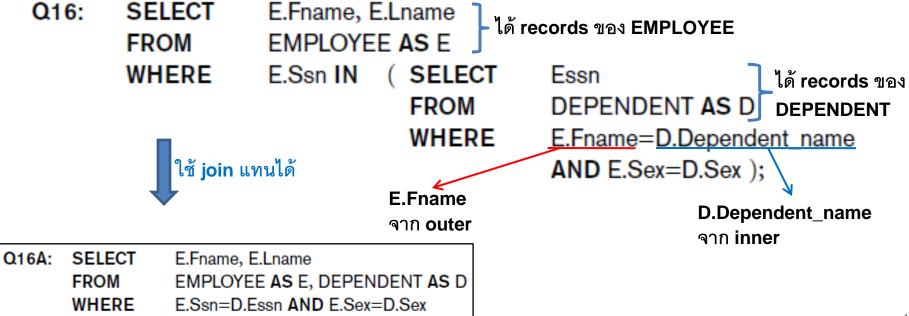
ถ้าใช้ Salary > SOME หรือ
Salary > ANY
ผลลัพธ์คือ พนักงานที่มีเงินเดือน
มากกว่าบางคนที่อยู่ในแผนก
หมายเลข 5

Avoid potential errors and ambiguities

AND E.Fname=D.Dependent_name;

 Create tuple variables (aliases) for all tables referenced in SQL query

Query 16. Retrieve the name of each employee who has a dependent with the same first name and is the same sex as the employee.



Correlated Nested Queries

- Correlated nested query
 - Evaluated once for each tuple in the outer query
- EXISTS function
 - Check whether the result of a correlated nested query is empty or not
 - ตรวจสอบ outer กับ inner ถ้าเจอใน inner แค่ 1 ก็ return True ไม่ต้องเช็คให้ถึง record สุดท้ายของ inner

 http://www.dba-oracle.com/t_in_vs_exists_sql.htm

SELECT *
FROM customers
WHERE EXISTS (SELECT *

The Oracle documentation notes that:

"If the selective predicate is in the subquery, then use IN. If the selective predicate is in the parent query, then use EXISTS."

FROM order_details
WHERE customers.customer_id = order_details.customer_id);

→ ลูกค้าที่มี order ใน order_details อะไรก็ได้ (ขอให้มีอย่างน้อย 1 order)

The EXISTS and UNIQUE Functions in SQL

- EXISTS and NOT EXISTS
 - Typically used in conjunction with a correlated nested query

```
SELECT *
FROM Customers
WHERE NOT EXISTS (SELECT *
FROM order_details
WHERE customers.customer_id = order_details.customer_id);
→ ลูกค้าที่ไม่มี order ใน order_details
```

- SQL function UNIQUE(Q)
 - Returns TRUE if there are no duplicate tuples in the result of query Q

```
SELECT *
FROM Student as S
WHERE Unique (SELECT CourseID
FROM Enroll as E
WHERE S.StudentID = E.StudentID);
```

→ นักศึกษาที่ลงทะเบียนแต่ละรายวิชาครั้งเดียว

Explicit Sets and Renaming of Attributes in SQL

- Can use explicit set of values in WHERE clause
- Use qualifier AS followed by desired new name
 - Rename any attribute that appears in the result of a query

Q8A: SELECT E.Lname AS Employee_name, S.Lname AS Supervisor_name

FROM EMPLOYEE AS E, EMPLOYEE AS S

WHERE E.Super_ssn=S.Ssn;

Joined Tables in SQL and Outer Joins

Joined table

 Permits users to specify a table resulting from a join operation in the FROM clause of a query

The FROM clause in Q1A

Contains a single joined table

Q1A: SELECT Fname, Lname, Address

FROM (EMPLOYEE JOIN DEPARTMENT ON Dno=Dnumber)

WHERE Dname='Research';

SELECT Customers.CustomerName, Orders.OrderID

FROM Customers INNER JOIN Orders ON Customers.CustomerID=Orders.CustomerID

ORDER BY Customers.CustomerName;

Joined Tables in SQL and Outer Joins (cont'd.)

- Specify different types of join
 - NATURAL JOIN
 - Various types of OUTER JOIN
- NATURAL JOIN on two relations R and S
 - No join condition specified
 - Implicit EQUIJOIN condition for each pair of attributes with same name from R and S

Joined Tables in SQL and Outer Joins (cont'd.)

Inner join

- Default type of join in a joined table
- Tuple is included in the result only if a matching tuple exists in the other relation
- LEFT OUTER JOIN
 - Every tuple in left table must appear in result
 - If no matching tuple
 - Padded with NULL values for attributes of right table
- RIGHT OUTER JOIN
 - Every tuple in right table must appear in result
 - If no matching tuple
 - Padded with NULL values for the attributes of left table
- FULL OUTER JOIN
- Can nest join specifications

Joined Tables in SQL and Outer Joins (cont'd.)

Reference http://en.wikipedia.org/wiki/Join_(SQL)

Employee table					
LastName	DepartmentID				
Rafferty	31				
Jones	33				
Heisenberg	33				
Robinson	34				
Smith	34				
Williams	NULL				

Department table				
DepartmentID	DepartmentName			
31	Sales			
33	Engineering			
34	Clerical			
35	Marketing			

SELEC	CT *					
FROM	employee	LEFT	OUTER	JOIN	department	
ON	employee	.Depar	rtment]	[D =	department.De	partmentID;

Employee.LastName	Employee.DepartmentID	Department.DepartmentName	Department.DepartmentID
Jones	33	Engineering	33
Rafferty	31	Sales	31
Robinson	34	Clerical	34
Smith	34	Clerical	34
Williams	NULL	NULL	NULL
Heisenberg	33	Engineering	33

SELECT *
FROM employee RIGHT OUTER JOIN department
ON employee.DepartmentID = department.DepartmentID;

Employee.LastName	Employee.DepartmentID	Department.DepartmentName	Department.DepartmentID
Smith	34	Clerical	34
Jones	33	Engineering	33
Robinson	34	Clerical	34
Heisenberg	33	Engineering	33
Rafferty	31	Sales	31
NULL	NULL	Marketing	35

FROM employee FULL OUTER JOIN department
ON employee.DepartmentID = department.DepartmentID;

Employee.LastName	Employee.DepartmentID	Department.DepartmentName	Department.DepartmentID
Smith	34	Clerical	34
Jones	33	Engineering	33
Robinson	34	Clerical	34
Williams	NULL	NULL	NULL
Heisenberg	33	Engineering	33
Rafferty	31	Sales	31
NULL	NULL	Marketing	35

More Complex SQL Insert, Update, Delete

More complex of insert command:

```
INSERT INTO first_table_name [(column1, column2, ..., columnN)]
    SELECT column1, column2, ..., columnN
    FROM second_table_name
    [WHERE condition];
```

INSERT INTO Customers (CustomerName, Country)
SELECT SupplierName, Country
FROM Suppliers;

```
INSERT INTO Customers (CustomerName, Country)
SELECT SupplierName, Country
FROM Suppliers
WHERE Country='Germany';
```

[]: optional

More Complex SQL Insert, Update, Delete

More complex of update command:

```
UPDATE ips
SET countryid = (select countryid from country where ips.iso=country.iso);
```

```
UPDATE country p, ips pp
SET pp.countryid = p.countryid
WHERE pp.iso = p.iso;
```

```
UPDATE [table1_name] AS t1 INNER JOIN [table2_name] AS t2 ON t1.[column1_name] = t2.[column1_name] SET t1.[column2_name] = t2.[column2_name];
```

```
UPDATE business AS b INNER JOIN business_geocode AS g ON b.business_id = g.business_id
SET b.mapx = g.latitude, b.mapy = g.longitude
WHERE (b.mapx = " or b.mapx = 0) and g.latitude > 0;
```

More Complex SQL Insert, Update, Delete

 Additional features allow users to specify more complex of delete command:

```
PELETE t1, t2
FROM t1 INNER JOIN t2 INNER JOIN t3
WHERE t1.id=t2.id AND t2.id=t3.id;
```

```
DELETE a1, a2
FROM db1.t1 AS a1 INNER JOIN db2.t2 AS a2
WHERE a1.id=a2.id;
```

```
DELETE w
FROM WorkRecord2 w INNER JOIN Employee e ON EmployeeRun = EmployeeNo
WHERE Company = '1' AND Date = '2013-05-06'
```

```
DELETE WorkRecord2, Employee
FROM WorkRecord2 INNER JOIN Employee ON (EmployeeRun = EmployeeNo)
WHERE Company = '1' AND Date = '2013-05-06';
```

Aggregate Functions in SQL

- Used to summarize information from multiple tuples into a single-tuple summary
- Grouping: ข้อมูลสรุปแยกตามกลุ่ม
 - Create subgroups of tuples before summarizing
- Built-in aggregate functions
 - COUNT, SUM, MAX, MIN, and AVG
- Functions can be used in the SELECT clause or in a HAVING clause

Aggregate Functions in SQL (cont'd.)

• NULL values discarded when aggregate functions are applied to a particular column: ค่า NULL จะไม่ถูกรวมเมื่อ aggregation function ใช้กับ column เช่น ถ้า Salary เป็น NULL และใช้ ฟังก์ชัน Average ค่า NULL จะไม่ถูกรวม

EMPLOYEE

Fname	Minit	Lname	Ssn	Bdate	Address	Sex	Salary	Super_ssn	Dno
John	В	Smith	123456789	1965-01-09	731 Fondren, Houston, TX	М	30000	333445555	5
Franklin	Т	Wong	333445555	1955-12-08	638 Voss, Houston, TX	М	40000	888665555	5
Alicia	J	Zelaya	999887777	1968-01-19	3321 Castle, Spring, TX	F	25000	987654321	4
Jennifer	S	Wallace	987654321	1941-06-20	291 Berry, Bellaire, TX	F	43000	888665555	4
Ramesh	K	Narayan	666884444	1962-09-15	975 Fire Oak, Humble, TX	М	38000	333445555	5
Joyce	Α	English	453453453	1972-07-31	5631 Rice, Houston, TX	F	25000	333445555	5
Ahmad	٧	Jabbar	987987987	1969-03-29	980 Dallas, Houston, TX	М	25000	987654321	4
James	Е	Borg	888665555	1937-11-10	450 Stone, Houston, TX	М	55000	NULL	1

DEPARTMENT

Dname	Dnumber	Mgr_ssn	Mgr_start_date
Research	5	333445555	1988-05-22
Administration	4	987654321	1995-01-01
Headquarters	1	888665555	1981-06-19

Q21: SELECT COUNT (*)
FROM EMPLOYEE;

Query 20. Find the sum of the salaries of all employees of the 'Research' department, as well as the maximum salary, the minimum salary, and the average salary in this department.

Q20: SELECT SUM (Salary), MAX (Salary), MIN (Salary), AVG (Salary)

FROM (EMPLOYEE JOIN DEPARTMENT ON Dno=Dnumber)

WHERE Dname='Research';

Q22: SELECT COUNT (*)

FROM EMPLOYEE, DEPARTMENT

WHERE DNO=DNUMBER AND DNAME='Research';

Queries 21 and 22. Retrieve the total number of employees in the company (Q21) and the number of employees in the 'Research' department (Q22).

Grouping: The GROUP BY and HAVING Clauses

- Partition relation into subsets of tuples
 - Based on grouping attribute(s)
 - Apply function to each such group independently
- GROUP BY clause
 - Specifies grouping attributes
- If NULLs exist in grouping attribute
 - Separate group created for all tuples with a NULL value in grouping attribute

Query 24. For each department, retrieve the department number, the number of employees in the department, and their average salary.

55000

Q24: SELECT Dno, COUNT (*), AVG (Salary)
FROM EMPLOYEE

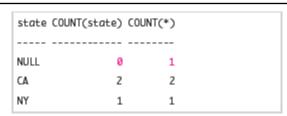
888665555

GROUP BY Dno;

Minit Lname Ssn Salary Count (*) Avg (Salary) Fname Super ssn Dno Dno В John Smith 123456789 30000 333445555 5 33250 Franklin Wong 888665555 5 333445555 40000 31000 666884444 333445555 Ramesh Narayan 38000 5 55000 Result of Q24 English 453453453 25000 333445555 5 Joyce Zelaya 987654321 Alicia 999887777 25000 Wallace 987654321 43000 888665555 Jennifer Ahmad ٧ Jabbar 987987987 25000 987654321

NULL

SELECT state, COUNT(state),
COUNT(*)
FROM publishers
GROUP BY state;



Bong

James

Grouping: The GROUP BY and HAVING Clauses (cont'd.)

HAVING clause

Provides a condition on the summary information

Query 26. For each project on which more than two employees work, retrieve the project number, the project name, and the number of employees who work on the project.

Q26: SELECT Pnumber, Pname, COUNT (*)

FROM PROJECT, WORKS_ON

WHERE Pnumber=Pno
GROUP BY Pnumber, Pname
HAVING COUNT (*) > 2;

Pname	<u>Pnumber</u>		Essn	Pno	Hours	_
ProductX	1		123456789	1	32.5	
ProductX	1		453453453	1	20.0	
ProductY	2		123456789	2	7.5	
ProductY	2		453453453	2	20.0	
ProductY	2		333445555	2	10.0	_
ProductZ	3		666884444	3	40.0	
ProductZ	3		333445555	3	10.0	_
Computerization	10		333445555	10	10.0	
Computerization	10		999887777	10	10.0	
Computerization	10		987987987	10	35.0	_
Reorganization	20	1	333445555	20	10.0	Ī
Reorganization	20		987654321	20	15.0	
Reorganization	20		888665555	20	NULL	_
Newbenefits	30		987987987	30	5.0	
Newbenefits	30		987654321	30	20.0	
Newbenefits	30	1	999887777	30	30.0	İ

 These groups are not selected by the HAVING condition of Q26.

Pname	<u>Pnumber</u>	 <u>Essn</u>	<u>Pno</u>	Hours
ProductY	2	123456789	2	7.5
ProductY	2	453453453	2	20.0
ProductY	2	333445555	2	10.0
Computerization	10	333445555	10	10.0
Computerization	10	 999887777	10	10.0
Computerization	10	987987987	10	35.0
Reorganization	20	333445555	20	10.0
Reorganization	20	987654321	20	15.0
Reorganization	20	888665555	20	NULL
Newbenefits	30	987987987	30	5.0
Newbenefits	30	987654321	30	20.0
Newbenefits	30	999887777	30	30.0

After applying the HAVING clause condition

Count (*)

3

3

3

Pname

ProductY

Computerization

Reorganization

Newbenefits
Result of Q26
(Pnumber not shown)



Grouping: The GROUP BY and HAVING Clauses (cont'd.)

Query 28. For each department that has more than five employees, retrieve the department number and the number of its employees who are making more than \$40,000.

Q28: SELECT Dnumber, COUNT (*)

FROM DEPARTMENT, EMPLOYEE

WHERE Dnumber=Dno AND Salary>40000 AND

(SELECT Dno

FROM EMPLOYEE

GROUP BY Dno

HAVING COUNT (*) > 5

Output =?

EMPLOYEE

Fname	Minit	Lname	Ssn	Bdate	Address	Sex	Salary	Super_ssn	Dno
John	В	Smith	123456789	1965-01-09	731 Fondren, Houston, TX	М	30000	333445555	5
Franklin	T	Wong	333445555	1955-12-08	638 Voss, Houston, TX	М	40000	888665555	5
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Ahmad	٧	Jabbar	987987987	1969-03-29	980 Dallas, Houston, TX	М	25000	987654321	4
James	Е	Borg	888665555	1937-11-10	450 Stone, Houston, TX	М	55000	NULL	1

DEPARTMENT

Dname	Dnumber	Mgr_ssn	Mgr_start_date
Research	5	333445555	1988-05-22
Administration	4	987654321	1995-01-01
Headquarters	1	888665555	1981-06-19

Discussion and Summary of SQL Queries

```
SELECT <attribute and function list>
FROM 
[ WHERE <condition> ]
[ GROUP BY <grouping attribute(s)> ]
[ HAVING <group condition> ]
[ ORDER BY <attribute list> ];
```

Specifying Constraints as Assertions and Actions as Triggers

CREATE ASSERTION

 Specify additional types of constraints outside scope of built-in relational model constraints

CREATE TRIGGER

Specify automatic actions that database system will perform when certain events and conditions occur



Specifying General Constraints as Assertions in SQL

- CREATE ASSERTION
 - Specify a query that selects any tuples that violate the desired condition
 - Use only in cases where it is not possible to use
 CHECK on attributes and domains

```
General Syntax: CREATE ASSERTION < name > CHECK(< condition >)
```

```
CREATE ASSERTION SALARY_CONSTRAINT
CHECK ( NOT EXISTS ( SELECT *
```

FROM EMPLOYEE E, EMPLOYEE M,

DEPARTMENT D

WHERE E.Salary>M.Salary

AND E.Dno=D.Dnumber

AND D.Mgr_ssn=M.Ssn));

Introduction to Triggers in SQL

- CREATE TRIGGER statement
 - Used to monitor the database
- Typical trigger has three components:
 - Event(s): e.g. insert, update
 - Condition
 - Action: a sequence of SQL statements

```
CREATE TRIGGER salary_trigger

BEFORE UPDATE ON employee_table REFERENCING NEW ROW AS n, OLD ROW AS o

FOR EACH ROW IF n.salary <> o.salary THEN
...

END IF; ;
```

Views (Virtual Tables) in SQL

- Concept of a view in SQL
 - Single table derived from other tables
 - Considered to be a virtual table
- CREATE VIEW command
 - Give table name, list of attribute names, and a query to specify the contents of the view

V1:	CREATE VIEW AS SELECT FROM WHERE	WORKS_ON1 Fname, Lname, Pname, Hours EMPLOYEE, PROJECT, WORKS_ON Ssn=Essn AND Pno=Pnumber;
V2:	CREATE VIEW AS SELECT FROM WHERE GROUP BY	DEPT_INFO(Dept_name, No_of_emps, Total_sal) Dname, COUNT (*), SUM (Salary) DEPARTMENT, EMPLOYEE Dnumber=Dno Dname;

Specification of Views in SQL (cont'd.)

- Specify SQL queries on a view
- View always up-to-date
 - Responsibility of the DBMS and not the user
- DROP VIEW command
 - Dispose of a view

View Implementation, View Update, and Inline Views

- Complex problem of efficiently implementing a view for querying
- Query modification approach
 - Modify view query into a query on underlying base tables
 - Disadvantage: inefficient for views defined via complex queries that are time-consuming to execute (กรณี join หลาย table จะใช้เวลานาน)

V1: CREATE VIEW WORKS_ON1

AS SELECT Fname, Lname, Pname, Hours

FROM EMPLOYEE, PROJECT, WORKS_ON WHERE Ssn=Essn AND Pno=Pnumber;

QV1: S

FROM WORKS_ON1
WHERE Pname='ProductX';



SELECT FROM WHERE Fname, Lname
EMPLOYEE, PROJECT, WORKS_ON
Ssn=Essn AND Pno=Pnumber
AND Pname='ProductX';

View Implementation

View materialization approach

- Physically create a temporary view table when the view is first queried
- Keep that table on the assumption that other queries on the view will follow
- Requires efficient strategy for automatically updating the view table when the base tables are updated

Incremental update strategies

 DBMS determines what new tuples must be inserted, deleted, or modified in a materialized view table

View Update and Inline Views

- Update on a view defined on a single table without any aggregate functions
 - Can be mapped to an update on underlying base table
- View involving joins
 - Often not possible for DBMS to determine which of the updates is intended

```
(a):
       UPDATEWORKS ON
       SET
                                    Pnumber
                  Pno =
                          SELECT
                           FROM
                                    PROJECT
                                    Pname='ProductY')
                          WHERE
       WHERE
                  Essn IN ( SELECT
                                    Ssn
                          FROM
                                    EMPLOYEE
                                    Lname='Smith' AND Fname='John'
                          WHERE
                  AND
                  Pno =
                          SELECT
                                    Pnumber
                                    PROJECT
                           FROM
                                    Pname='ProductX');
                          WHERE
                                Pname = 'ProductY'
(b):
       UPDATE PROJECT
                         SET
                  Pname = 'ProductX';
```

การ Update Base Table ที่เกี่ยวข้องกับ UV1 UV1 จะต้องถูก Update ด้วย ซึ่งอาจมีผลทำให้ View อื่นที่อ้างอิงถึง UV1 แต่ใช้ข้อมูลต่าง มุมมอง อาจจะได้ข้อมูลที่ไม่ถูกต้องได้

```
UV1: UPDATE WORKS_ON1
SET Pname = 'ProductY'
WHERE Lname='Smith' AND Fname='John'
AND Pname='ProductX';
```

View Update and Inline Views (cont'd.)

Clause WITH CHECK OPTION

 Must be added at the end of the view definition if a view is to be updated

```
CREATE VIEW VW_TechnicianEmployees
AS SELECT EmployeeID, Title, ManagerID
FROM HumanResources. Employee
WHERE <u>Title LIKE '%technician%'</u>
WITH CHECK OPTION;
```

```
UPDATE VW_TechnicianEmployees
SET Title = 'Chief'
WHERE EmployeeID=13
```

In-line view

- Defined in the FROM clause of an SQL query
- ใช้ลดความซับซ้อนของ Query

Schema Change Statements in SQL

Schema evolution commands

- Can be done while the database is operational
- Does not require recompilation of the database schema

DROP command

- Used to drop named schema elements, such as tables, domains, or constraint
- Drop behavior options:
 - CASCADE and RESTRICT

Example:

- DROP SCHEMA COMPANY CASCADE;
 - **CASCADE** Automatically drop objects (tables, functions, etc.) that are contained in the schema.
 - **RESTRICT** Refuse to drop the schema if it contains any objects. This is the default.

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The ALTER Command

Alter table actions include:

- Adding or dropping a column (attribute)
- Changing a column definition
- Adding or dropping table constraints

Example:

- -ALTER TABLE COMPANY.EMPLOYEE ADD COLUMN Job VARCHAR(12);
- To drop a column
 - Choose either CASCADE or RESTRICT

The ALTER Command (cont'd.)

- Change constraints specified on a table
 - Add or drop a named constraint

ALTER TABLE COMPANY.EMPLOYEE

DROP CONSTRAINT EMPSUPERFK CASCADE;

Summary

- Complex SQL:
 - Nested queries, joined tables, outer joins, aggregate functions, grouping
- CREATE ASSERTION and CREATE TRIGGER
- Views
 - Virtual or derived tables